

Orchestrating Decision Engines

University of Verona, 2011

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Satisfiability Modulo Theories (SMT)

A Satisfiability Checker
with built-in support for useful theories

Satisfiability Modulo Theories (SMT)

$$b + 2 = c \text{ and } f(\text{read}(\text{write}(a, b, 3), c-2) \neq f(c-b+1)$$

Satisfiability Modulo Theories (SMT)

$$b + 2 = c \text{ and } f(\text{read}(\text{write}(a, b, 3), c-2) \neq f(c-b+1)$$

Arithmetic

Satisfiability Modulo Theories (SMT)

$b + 2 = c$ and $f(\text{read}(\text{write}(a, b, 3), c-2) \neq f(c-b+1)$

Array Theory

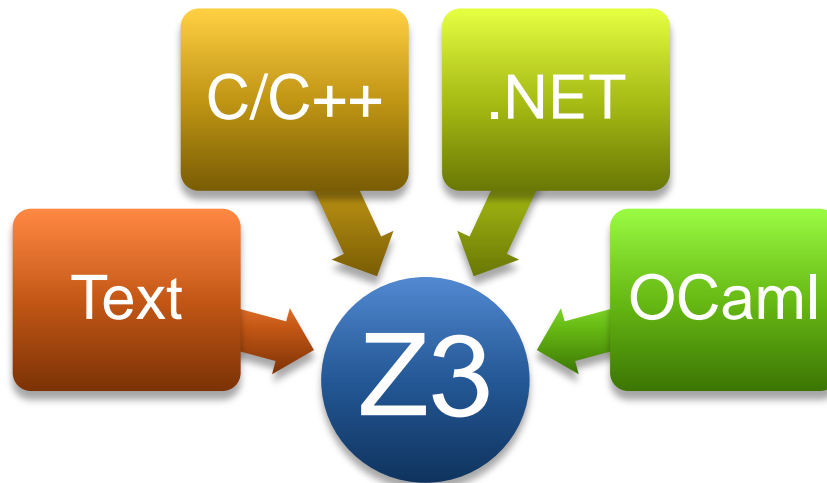
Satisfiability Modulo Theories (SMT)

$$b + 2 = c \text{ and } f(\text{read}(\text{write}(a, b, 3), c - 2) \neq f(c - b + 1)$$

Uninterpreted
Functions

SMT@Microsoft: Solver

- Z3 is a solver developed at Microsoft Research.
- Development/Research driven by internal customers.
- Free for non-commercial use.
- Interfaces:



- <http://research.microsoft.com/projects/z3>

Some Applications @ Microsoft



The Spec#
Programming System

HAVOC



Hyper-V

Microsoft | Virtualization 

Terminator T-2

VCC

NModel

SLAM
`if(!node->x) i ++ visprocs; end() *node){`



Vigilante

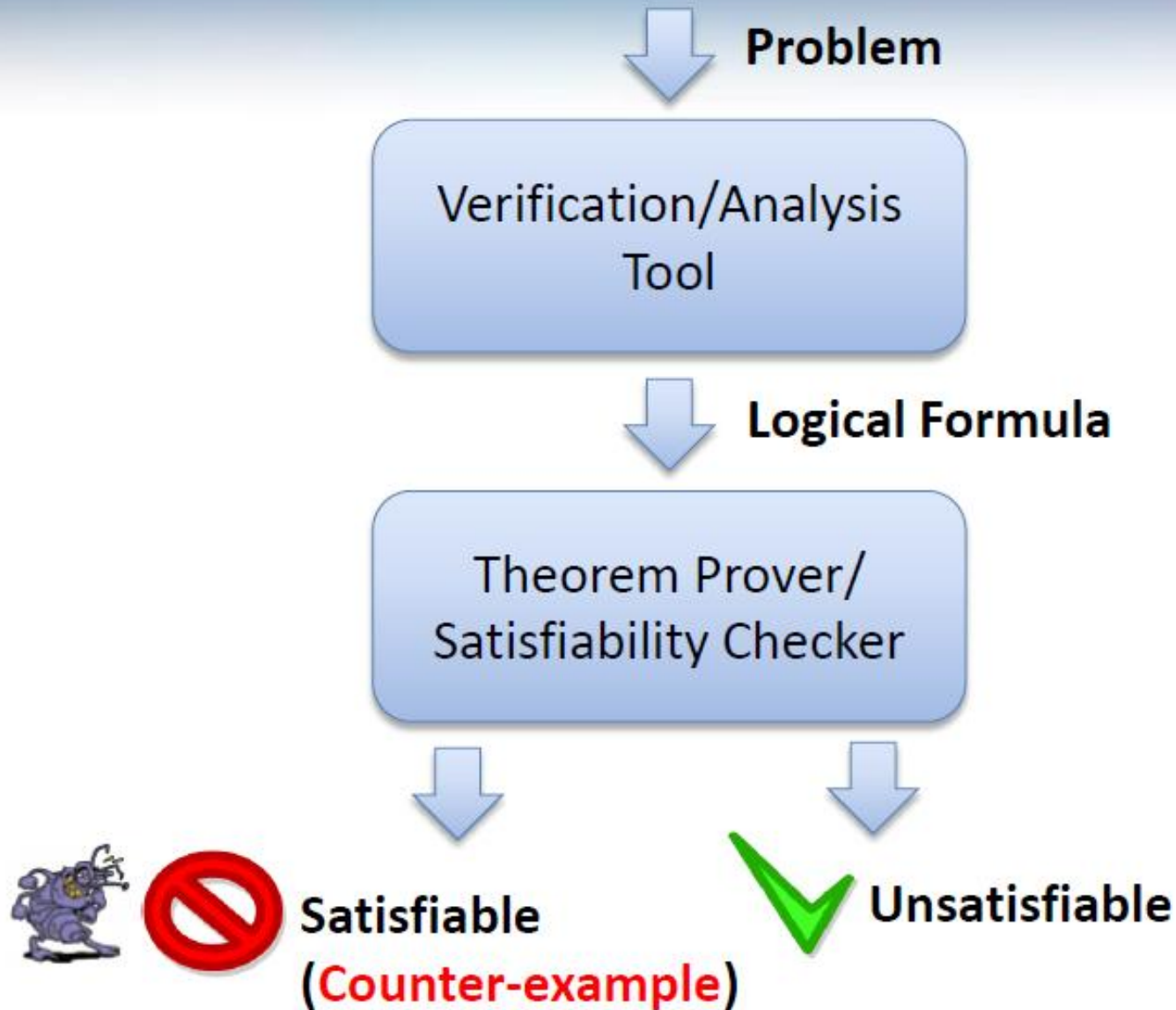
SpecExplorer



F7

SAGE

Verification/Analysis Tool: “Template”

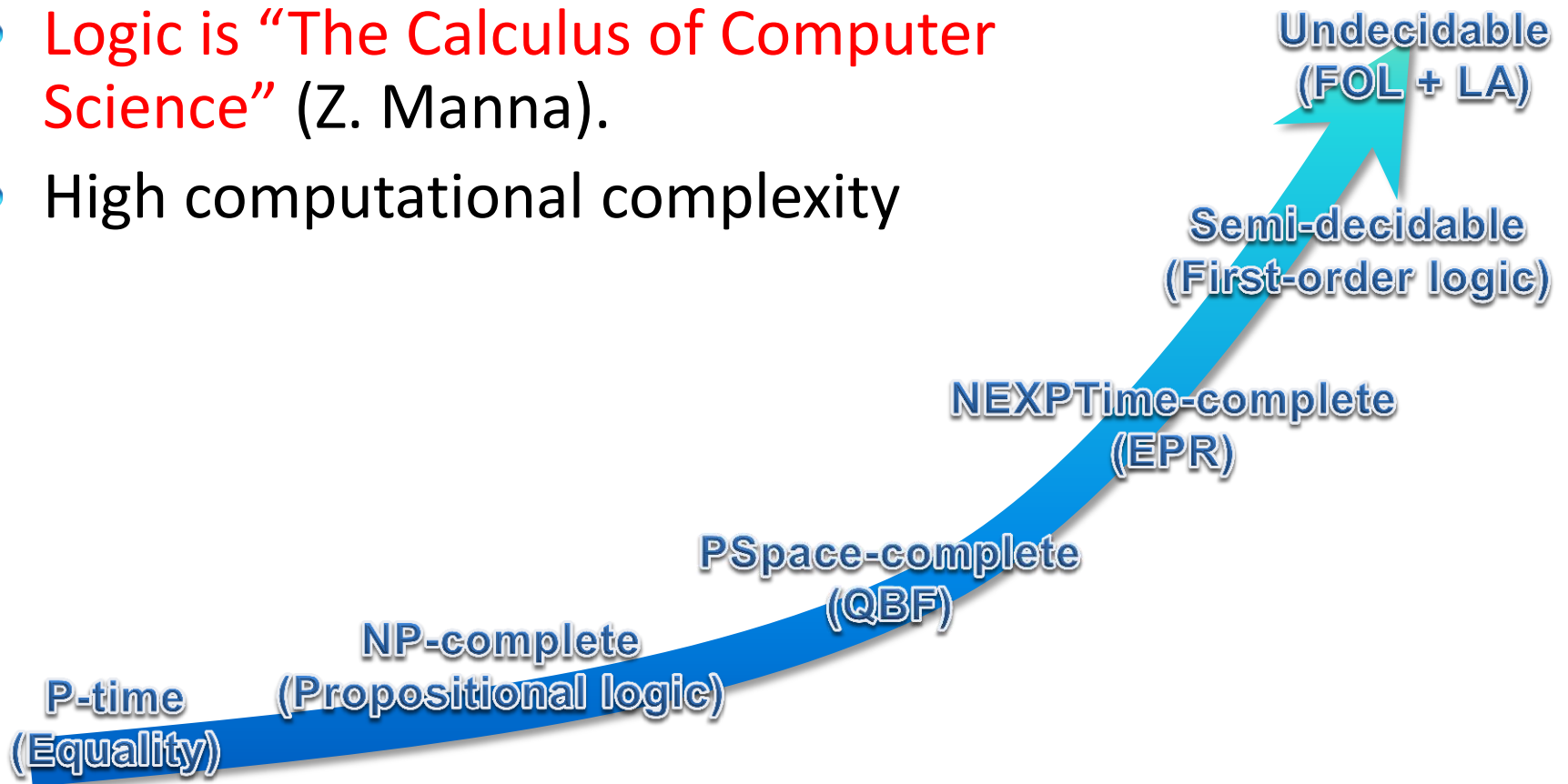


Symbolic Reasoning

Verification/Analysis tools
need some form of
Symbolic Reasoning

Symbolic Reasoning

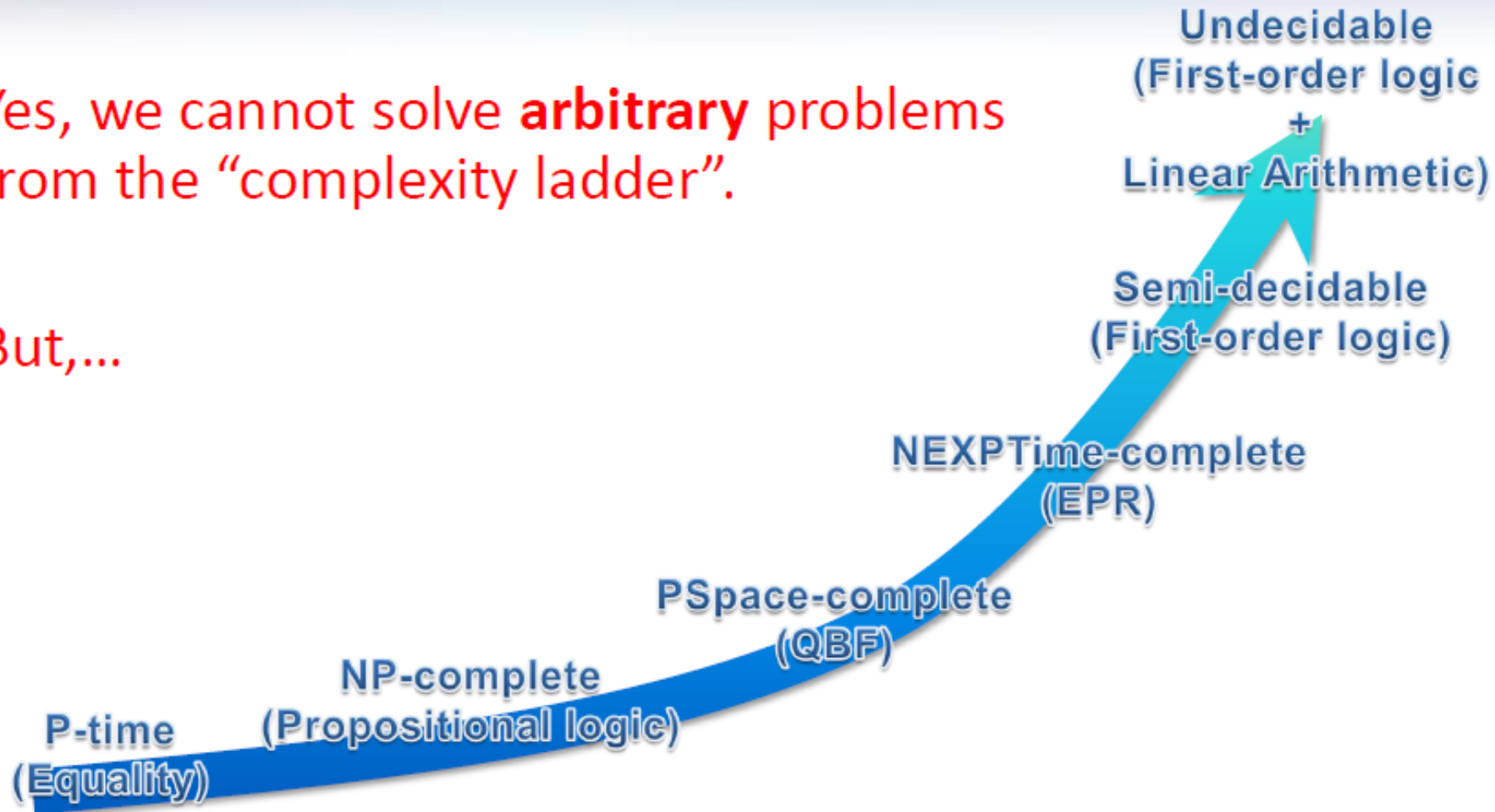
- Logic is “The Calculus of Computer Science” (Z. Manna).
- High computational complexity



Symbolic Reasoning

Yes, we cannot solve **arbitrary** problems from the “complexity ladder”.

But,...



Symbolic Reasoning

We can try to solve the
problems we find in
real applications

Applications

Test case generation

Verifying Compilers

Predicate Abstraction

Invariant Generation

Type Checking

Model Based Testing

Test-case generation

- Test (correctness + usability) is 95% of the deal:
 - Dev/Test is 1-1 in products.
 - Developers are responsible for unit tests.
- Tools:
 - Annotations and static analysis (SAL + ESP)
 - File Fuzzing
 - Unit test case generation

Security is critical

- Security bugs can be very expensive:
 - Cost of each MS Security Bulletin: \$600k to \$Millions.
 - Cost due to worms: \$Billions.
 - **The real victim is the customer.**
- Most security exploits are initiated via files or packets.
 - Ex: Internet Explorer parses dozens of file formats.
- Security testing: **hunting for million dollar bugs**
 - Write A/V
 - Read A/V
 - Null pointer dereference
 - Division by zero

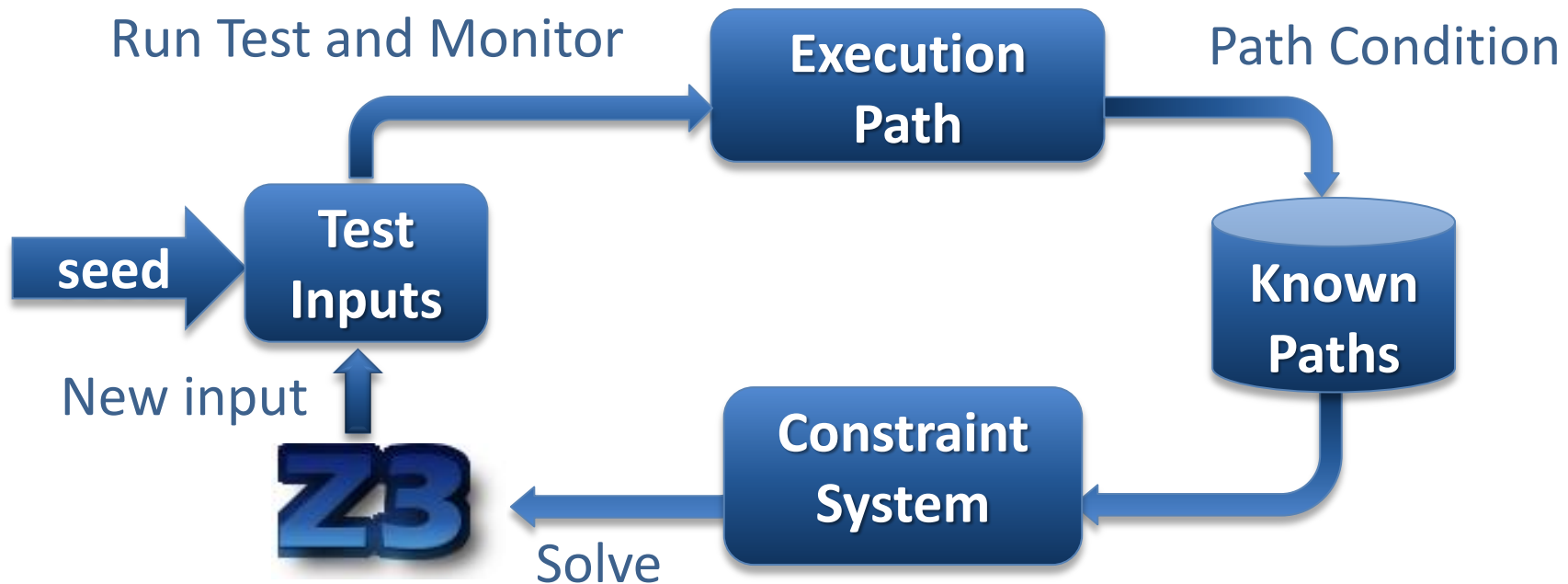


Hunting for Security Bugs.

- Two main techniques used by “*black hats*”:
 - Code inspection (of binaries).
 - *Black box fuzz testing.*
- **Black box** fuzz testing:
 - A form of black box random testing.
 - Randomly *fuzz* (=modify) a well formed input.
 - Grammar-based fuzzing: rules to encode how to fuzz.
- **Heavily** used in security testing
 - At MS: several internal tools.
 - Conceptually simple yet effective in practice



Directed Automated Random Testing (DART)



DARTish projects at Microsoft

PEX

Implements DART for .NET.

SAGE

Implements DART for x86 binaries.

YOGI

Implements DART to check the feasibility of program paths generated statically using a SLAM-like tool.

Vigilante

Partially implements DART to dynamically generate worm filters.

What is *Pex*?

- Test input generator
 - Pex starts from parameterized unit tests
 - Generated tests are emitted as traditional unit tests

ArrayList: The Spec

The image shows a screenshot of the MSDN .NET Framework Developer Center page for the `ArrayList.Add` method. The page has a red header with the MSDN logo and navigation tabs: Home, Library, Learn, Downloads, and Support. Below the header, there are links for 'Printer Friendly Version', 'Add To Favorites', 'Send', and 'Add Content...'. The main content area is titled '.NET Framework Class Library' and 'ArrayList.Add Method'. It describes the method's purpose: 'Adds an object to the end of the `ArrayList`.' It also specifies the 'Namespace: `System.Collections`' and 'Assembly: mscorlib (in mscorlib.dll)'. A 'Remarks' section explains that `ArrayList` accepts a null reference (Nothing in Visual Basic) as a valid value and allows duplicate elements. It further details that if the `Count` already equals the `Capacity`, the capacity is increased by automatically reallocating the internal array, and the existing elements are copied to the new array before the new element is added. If `Count` is less than `Capacity`, the method is an $O(1)$ operation. If the capacity needs to be increased to accommodate the new element, the method becomes an $O(n)$ operation, where n is `Count`. A 'Click to Rate and Give Feedback' link with a star rating is also visible.

msdn
.NET Framework Developer Center

Home Library Learn Downloads Support

Printer Friendly Version + Add To Favorites Send Add Content...

.NET Framework Class Library
ArrayList.Add Method

Adds an object to the end of the `ArrayList`.

Namespace: `System.Collections`
Assembly: mscorlib (in mscorlib.dll)

Remarks

`ArrayList` accepts a null reference (**Nothing** in Visual Basic) as a valid value and allows duplicate elements.

If `Count` already equals `Capacity`, the capacity of the `ArrayList` is increased by automatically reallocating the internal array, and the existing elements are copied to the new array before the new element is added.

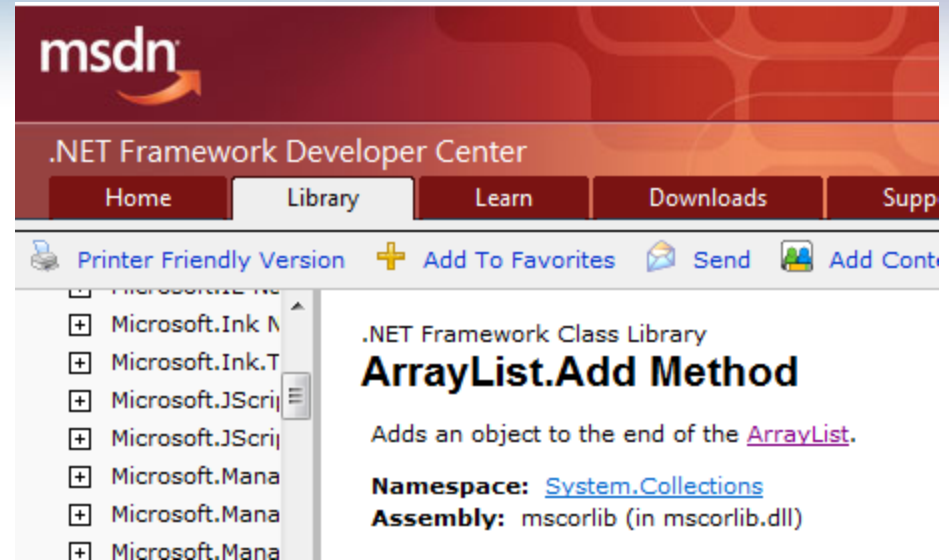
If `Count` is less than `Capacity`, this method is an $O(1)$ operation. If the capacity needs to be increased to accommodate the new element, this method becomes an $O(n)$ operation, where n is `Count`.

Click to Rate and Give Feedback ★★☆☆

ArrayList: AddItem Test

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...  
}
```



ArrayList: Starting Pex...

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...  
}
```

Inputs

ArrayList: Run 1, (0,null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...  
}
```

Inputs
(0,null)

ArrayList: Run 1, (0,null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...  
}
```

`c < 0` → false

Inputs	Observed Constraints
(0,null)	!(c<0)

ArrayList: Run 1, (0,null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item);  
    }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item;  
    }  
    ...  
}
```

0 == c → true

Inputs	Observed Constraints
(0,null)	!(c<0) && 0==c

ArrayList: Run 1, (0,null)

```
class ArrayListTest {  
  [PexMethod]  
  void AddItem(int c, object item) {  
    var list = new ArrayList(c);  
    list.Add(item);  
    Assert(list[0] == item); }  
}
```

item == item → true

```
class ArrayList {  
  object[] items;  
  int count;  
  
  ArrayList(int capacity) {  
    if (capacity < 0) throw ...;  
    items = new object[capacity];  
  }  
  
  void Add(object item) {  
    if (count == items.Length)  
      ResizeArray();  
  
    items[this.count++] = item; }  
  ...  
}
```

Inputs	Observed Constraints
(0,null)	!(c<0) && 0==c

ArrayList: Picking the next branch to cover

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...  
}
```

Constraints to
solve

Inputs

Observed
Constraints

(0,null)

!(c<0) && 0==c

!(c<0) && 0!=c



Z3

ArrayList: Solve constraints using SMT solver

```
class ArrayListTest {  
  [PexMethod]  
  void AddItem(int c, object item) {  
    var list = new ArrayList(c);  
    list.Add(item);  
    Assert(list[0] == item); }  
}
```

```
class ArrayList {  
  object[] items;  
  int count;  
  
  ArrayList(int capacity) {  
    if (capacity < 0) throw ...;  
    items = new object[capacity];  
  }  
  
  void Add(object item) {  
    if (count == items.Length)  
      ResizeArray();  
  
    items[this.count++] = item; }  
  ...  
}
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	



ArrayList: Run 2, (1, null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length) ResizeArray();  
  
        items[this.count++] = item; }  
    ...  
}
```

$0 == c \rightarrow \text{false}$

Constraints to solve	Inputs	Observed Constraints
	(0,null)	$!(c < 0) \ \&\& \ 0 == c$
$!(c < 0) \ \&\& \ 0 != c$	(1,null)	$!(c < 0) \ \&\& \ 0 != c$

ArrayList: Pick new branch

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...  
}
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	!(c<0) && 0!=c
c<0		



Z3

ArrayList: Run 3, (-1, null)

```
class ArrayListTest {  
  [PexMethod]  
  void AddItem(int c, object item) {  
    var list = new ArrayList(c);  
    list.Add(item);  
    Assert(list[0] == item); }  
}
```

```
class ArrayList {  
  object[] items;  
  int count;  
  
  ArrayList(int capacity) {  
    if (capacity < 0) throw ...;  
    items = new object[capacity];  
  }  
  
  void Add(object item) {  
    if (count == items.Length)  
      ResizeArray();  
  
    items[this.count++] = item; }  
  ...  
}
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	!(c<0) && 0!=c
c<0	(-1,null)	



Z3

ArrayList: Run 3, (-1, null)

```
class ArrayListTest {  
  [PexMethod]  
  void AddItem(int c, object item) {  
    var list = new ArrayList(c);  
    list.Add(item);  
    Assert(list[0] == item); }  
}
```

```
class ArrayList {  
  object[] items;  
  int count;  
  
  ArrayList(int capacity) {  
    if (capacity < 0) throw ...;  
    items = new object[capacity];  
  }  
  
  void Add(object item) {  
    if (count == items.Length)  
      ResizeArray();  
  
    items[this.count++] = item; }  
  ...  
}
```

$c < 0 \rightarrow \text{true}$

Constraints to solve	Inputs	Observed Constraints
	(0,null)	$!(c < 0) \ \&\& \ 0 == c$
$!(c < 0) \ \&\& \ 0 != c$	(1,null)	$!(c < 0) \ \&\& \ 0 != c$
$c < 0$	(-1,null)	$c < 0$

ArrayList: Run 3, (-1, null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...  
}
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	!(c<0) && 0!=c
c<0	(-1,null)	c<0

ArrayList: Run 3, (-1, null)

```
class ArrayListTest {  
    [PexMethod]  
    void AddItem(int c, object item) {  
        var list = new ArrayList(c);  
        list.Add(item);  
        Assert(list[0] == item); }  
}
```

```
class ArrayList {  
    object[] items;  
    int count;  
  
    ArrayList(int capacity) {  
        if (capacity < 0) throw ...;  
        items = new object[capacity];  
    }  
  
    void Add(object item) {  
        if (count == items.Length)  
            ResizeArray();  
  
        items[this.count++] = item; }  
    ...  
}
```

Constraints to solve	Inputs	Observed Constraints
	(0,null)	!(c<0) && 0==c
!(c<0) && 0!=c	(1,null)	!(c<0) && 0!=c
c<0	(-1,null)	c<0



White box testing in practice

How to test this code?

(Real code from .NET base class libraries.)

```
[SecurityPermissionAttribute(SecurityAction.LinkDemand, Flags=SecurityPermissionFlag.SerializationFormatter)]
public ResourceReader(Stream stream)
{
    if (stream==null)
        throw new ArgumentNullException("stream");
    if (!stream.CanRead)
        throw new ArgumentException(Environment.GetResourceString("Argument_StreamNotReadable"));

    _resCache = new Dictionary<String, ResourceLocator>(FastResourceComparer.Default);
    _store = new BinaryReader(stream, Encoding.UTF8);
    // We have a faster code path for reading resource files from an assembly.
    _ums = stream as UnmanagedMemoryStream;

    BCLDebug.Log("RESMGRFILEFORMAT", "ResourceReader .ctor(Stream). UnmanagedMemoryStream: "+(_ums!=null));
    ReadResources();
}
```

White box testing in practice

```
// Reads in the header information for a .resources file. Verifies some
// of the assumptions about this resource set, and builds the class table
// for the default resource file format.
private void ReadResources() {
    BCLDebug.Assert(_store != null, "ResourceReader is closed!");
    BinaryFormatter bf = new BinaryFormatter(null, new StreamingContext(StreamingContextStates.File |
#if !FEATURE_PAL
    _typeLimitingBinder = new TypeLimitingDeserializationBinder();
    bf.Binder = _typeLimitingBinder;
#endif

    _objFormatter = bf;
    try {
        // Read ResourceManager header
        // Check for magic number
        int magicNum = _store.ReadInt32();
        if (magicNum != ResourceManager.MagicNumber)
            throw new ArgumentException(Environment.GetResourceString("Resources_StreamNotValid"));
        // Assuming this is ResourceManager header v1 or greater, hopefully,
        // after the version number there is a number of bytes to skip
        // to bypass the rest of the ResMgr header.
        int resMgrHeaderVersion = _store.ReadInt32();
        if (resMgrHeaderVersion > 1) {
            int numBytesToSkip = _store.ReadInt32();
            BCLDebug.Log("RESMGRFILEFORMAT", LogLevel.Status, "ReadResources: Unexpected ResMgr header");
            BCLDebug.Assert(numBytesToSkip >= 0, "numBytesToSkip in ResMgr header should be positive!");
            _store.BaseStream.Seek(numBytesToSkip, SeekOrigin.Current);
        } else {
            BCLDebug.Log("RESMGRFILEFORMAT", "ReadResources: Parsing ResMgr header v1.");
            SkipInt32(); // We don't care about numBytesToSkip.
            // Read in type name for a suitable ResourceReader
            // ...
        }
    }
}
```

White box testing in practice

```
// Reads in the header information for a .resources file. Verifies some
// of the assumptions about this resource set, and builds the class table
// for the default resource file format.
private void ReadResources() {
    BCLDebug.Assert(_store != null, "ResourceReader is closed!");
    BinaryFormatter bf = new BinaryFormatter(null, new StreamingContext(StreamingContextStates.File |
#if !FEATURE_PAL
    _typeLimitingBinder = new TypeLimitingDeserializationBinder();
    bf.Binder = _typeLimitingBinder;
#endif

    _objFormatter = bf;
    try {
        // Read ResourceManager header
        // Check for magic number
        int magicNum = _store.ReadInt32();
        if public virtual int ReadInt32() {
            if (m_isMemoryStream) {
                // Read directly from MemoryStream buffer
                MemoryStream mStream = m_stream as MemoryStream;
                BCLDebug.Assert(mStream != null, "m_stream as MemoryStream != null");
                if
                    return mStream.InternalReadInt32();
            }
            else
            {
                FillBuffer(4);
                return (int) (m_buffer[0] | m_buffer[1] << 8 | m_buffer[2] << 16 | m_buffer[3] << 24);
            }
        }
    }
    // Read in type name for a suitable ResourceReader
    // ...
```

Pex – Test Input Generation

The image shows a screenshot of Microsoft Visual Studio with a project named 'TestProject1'. The main window displays the code for 'ResourceReaderTest1.cs'. The code is a C# class 'ResourceReaderTests' with a 'ParameterizedTest(byte[] a)' method. The method body includes a 'PexAssume.IsNotNull(a);' check, a 'fixed (byte* p = a)' statement, and a 'using (stream = new UnmanagedMemoryStream(p, a.Length))' block. Inside the using block, a 'ResourceReader' is instantiated and 'readEntries' is called. A red circle highlights the 'byte[] a' parameter in the method signature. A red arrow points from this circle to a callout box on the right. The callout box, titled 'Test input, generated by Pex', contains a C# code snippet: 'byte[] a = new byte[14];' followed by an array of values: 'a[0] = 206;', 'a[1] = 202;', 'a[2] = 239;', 'a[3] = 190;', 'a[7] = 64;', and 'a[11] = 100;'. A red circle highlights the 'ParameterizedTest(a);' line in this snippet. A blue arrow points from the 'ParameterizedTest(a);' line back to the 'ParameterizedTest' method signature in the main code. A context menu is open over the code, showing options like 'Pex it Ctrl + F8', 'Pex', 'Refactor', 'Insert Snippet...', 'Surround With...', 'Go To Definition', 'Find All References', 'Breakpoint', 'Run To Cursor', 'Cut', 'Copy', 'Paste', and 'Outlining'.

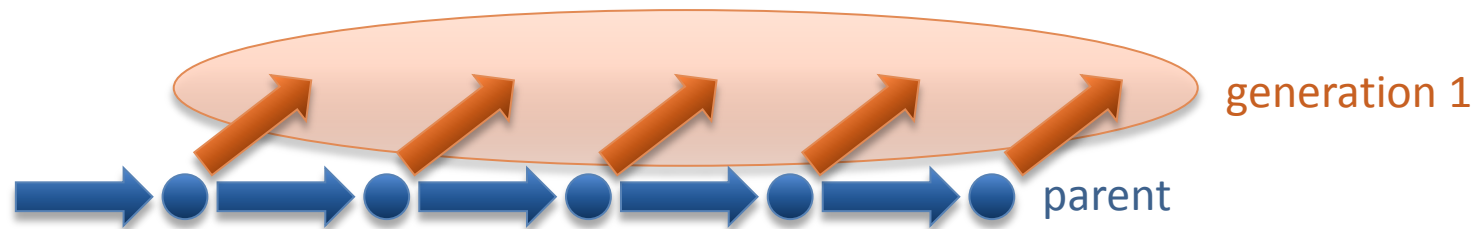
```
public class ResourceReaderTests
{
    [PexTest]
    public unsafe void ParameterizedTest(byte[] a)
    {
        PexAssume.IsNotNull(a);
        fixed (byte* p = a)
        using (stream = new UnmanagedMemoryStream(p, a.Length))
        {
            var reader = new ResourceReader(stream);
            readEntries(reader);
        }
    }
}
```

Test input, generated by Pex

```
byte[] a = new byte[14];
a[0] = 206;
a[1] = 202;
a[2] = 239;
a[3] = 190;
a[7] = 64;
a[11] = 100;
ParameterizedTest(a);
```


SAGE

- Apply DART to large applications (not units).
- Start with well-formed input (not random).
- Combine with generational search (not DFS).
 - Negate 1-by-1 each constraint in a path constraint.
 - Generate many children for each parent run.



Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 0 – seed file

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 00 00 00 00 00 00 00 00 00 00 00 00 ; RIFE.....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 1

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF....***.....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 2

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 3

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 00 00 00 00 ; ....trh.....
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 4

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh... vids
00000040h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 5

Zero to Crash in 10 Generations


- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 00 00 00 00 ; ....strf.....
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 6

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 28 00 00 00 ; ....strf.......
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 7

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 28 00 00 00 ; ....strf....(...)
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 C9 9D E4 4E ; .....EANI
00000060h: 00 00 00 00 ; ....
```

Generation 8

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 00 00 00 00 28 00 00 00 ; ....strf....(...)
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 01 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 9

Zero to Crash in 10 Generations

- Starting with 100 zero bytes ...
- SAGE generates a crashing test for Media1 parser

```
00000000h: 52 49 46 46 3D 00 00 00 00 ** ** ** 20 00 00 00 00 ; RIFF=...*** ....
00000010h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000020h: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 ; .....
00000030h: 00 00 00 00 73 74 72 68 00 00 00 00 76 69 64 73 ; ....strh....vids
00000040h: 00 00 00 00 73 74 72 66 B2 75 76 3A 28 00 00 00 ; ....stri2uv:(...
00000050h: 00 00 00 00 00 00 00 00 00 00 00 00 00 01 00 00 ; .....
00000060h: 00 00 00 00 ; ....
```

Generation 10 – CRASH

SAGE (cont.)

- SAGE is very effective at finding bugs.
- Works on large applications.
- Fully automated
- Easy to deploy (x86 analysis – any language)
- Used in various groups inside Microsoft
- Powered by Z3.

SAGE \leftrightarrow Z3

- Formulas are usually big conjunctions.
- SAGE uses only the bitvector and array theories.
- Pre-processing step has a huge performance impact.
 - Eliminate variables.
 - Simplify formulas.
- Early unsat detection.

Application Scenarios

“Big” and hard formulas

Thousands of “small” and easy formulas

Short timeout (< 5 secs)

Application Scenarios

“Big” and hard formulas



The
Spec#

Programming System

HAVOC

VCC

Thousands of “small” and easy formulas

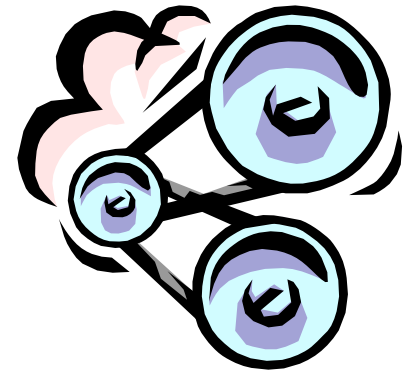
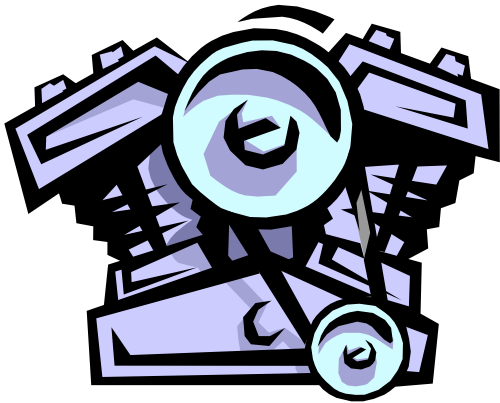


Short timeout (< 5secs)

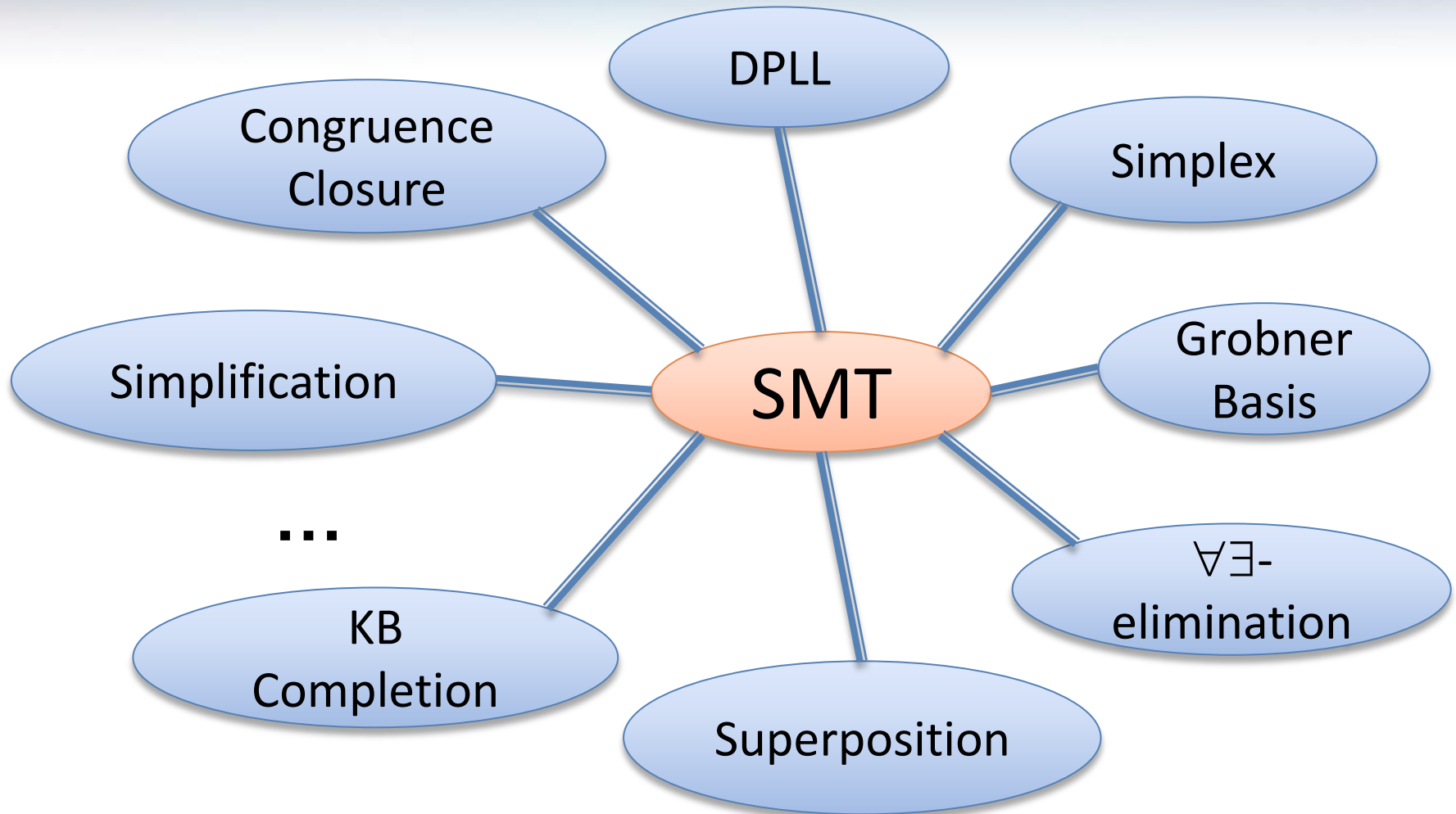
SAGE

Combining Engines

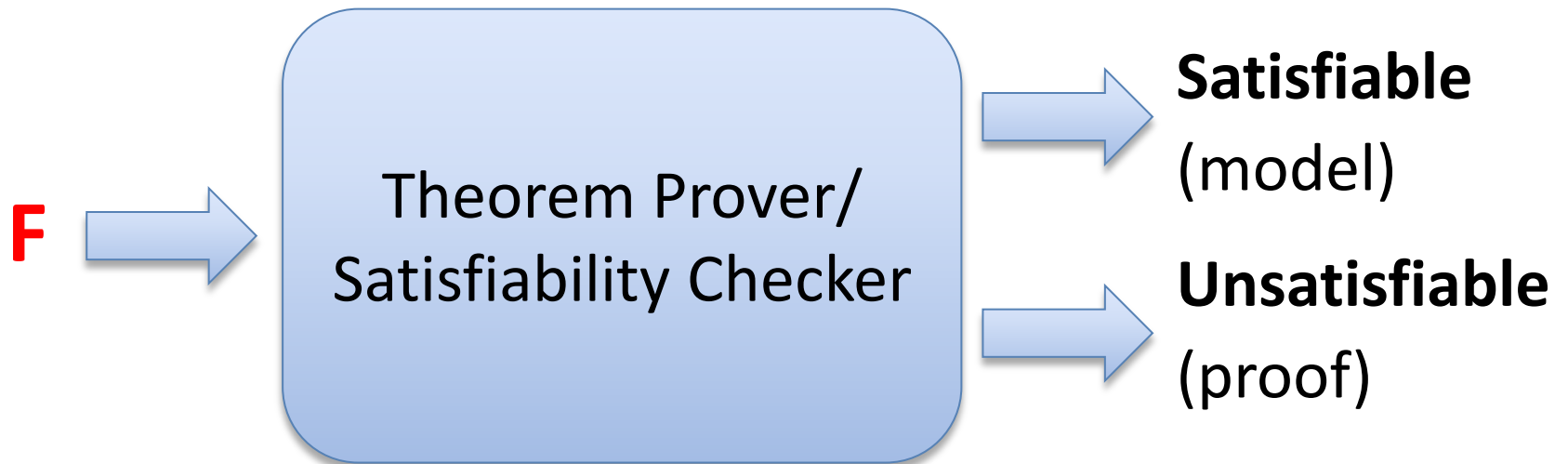
Current SMT solvers provide
a combination
of different engines



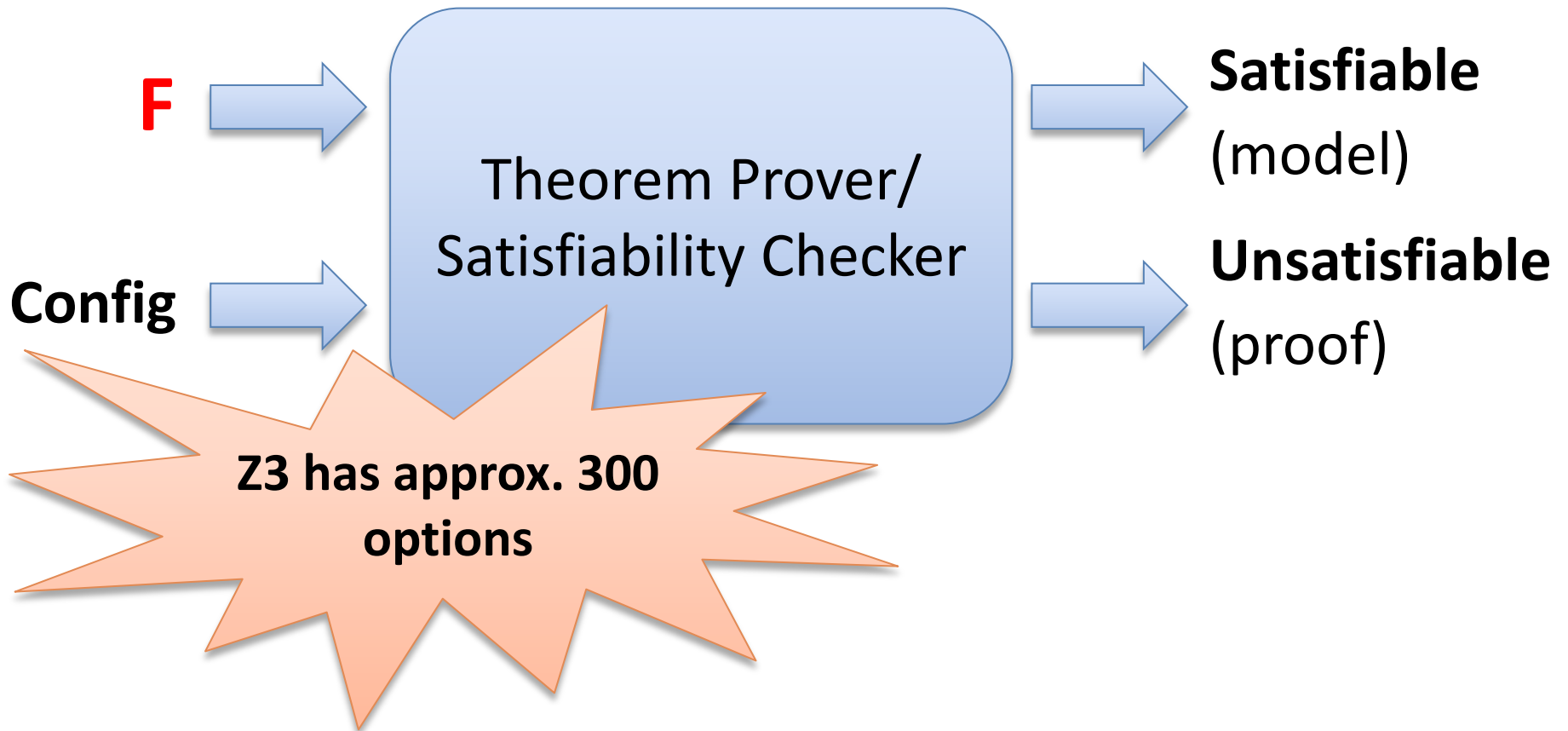
Combining Engines



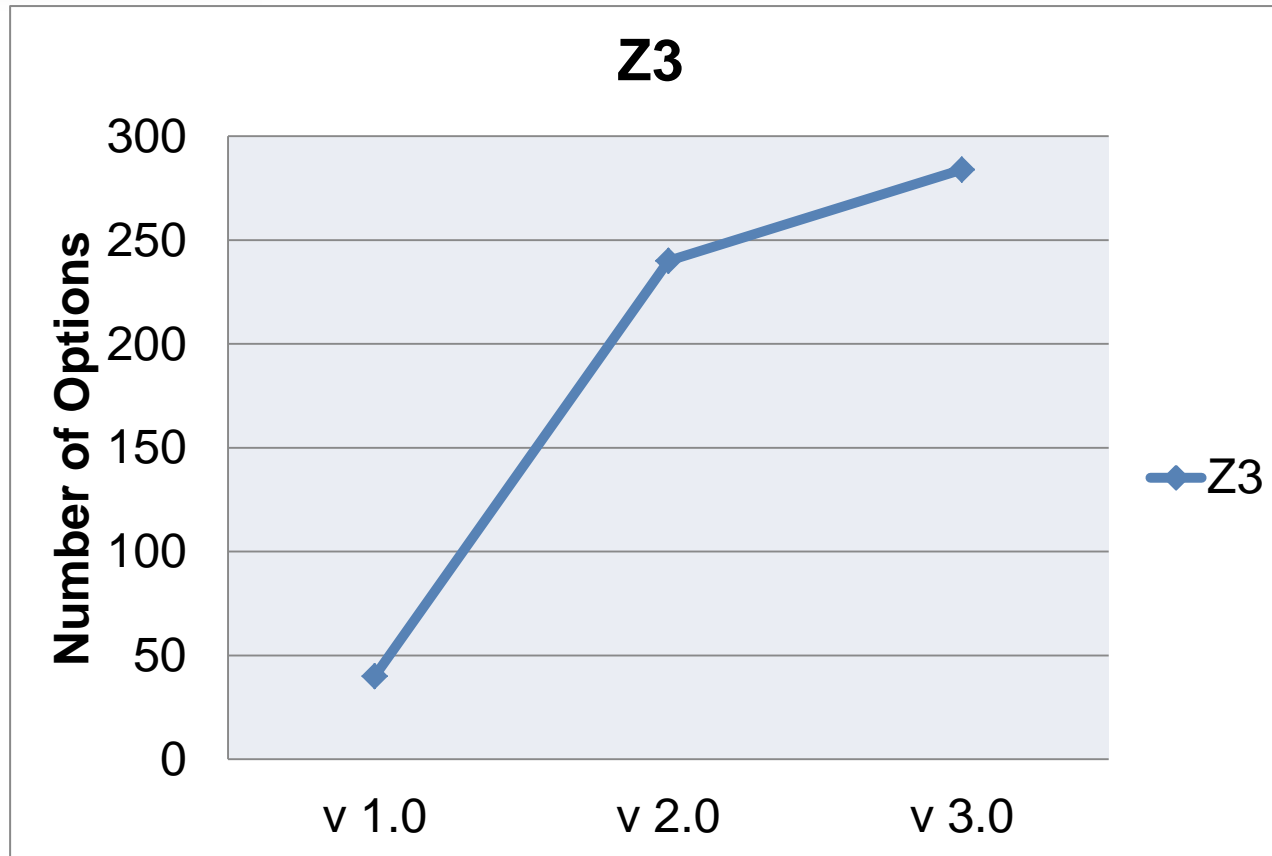
Theorem Provers & Satisfiability Checkers



Theorem Provers & Satisfiability Checkers



Z3 number of options evolution



Opening the “Black Box”

Actual feedback provided by Z3 users:

“Could you send me your CNF converter?”

“I want to implement my own search strategy.”

“I want to include these rewriting rules in Z3.”

“I want to apply a substitution to term t .”

“I want to compute the set of implied equalities.”

The Strategy Challenge

To build theoretical and practical tools allowing users to exert strategic control over core heuristic aspects of high performance SMT solvers.

What is a strategy?

Theorem proving as an exercise of combinatorial search

Strategies are **adaptations** of general search mechanisms which **reduce** the **search space** by tailoring its exploration to a **particular class** of formulas.

The Need for “Strategies”

Different Strategies for Different Domains.

The Need for “Strategies”

Different Strategies for Different Domains.

From timeout to 0.05 secs...

Example in Quantified Bit-Vector Logic (QBFV)

Join work with C. Wintersteiger and Y. Hamadi
FMCAD 2010

QBFV = Quantifiers + Bit-vectors + uninterpreted functions

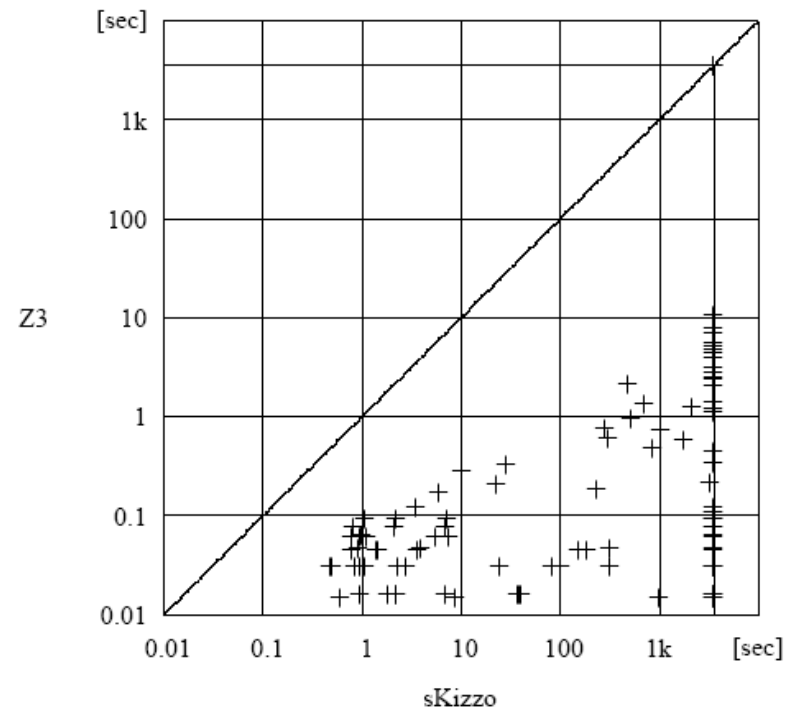
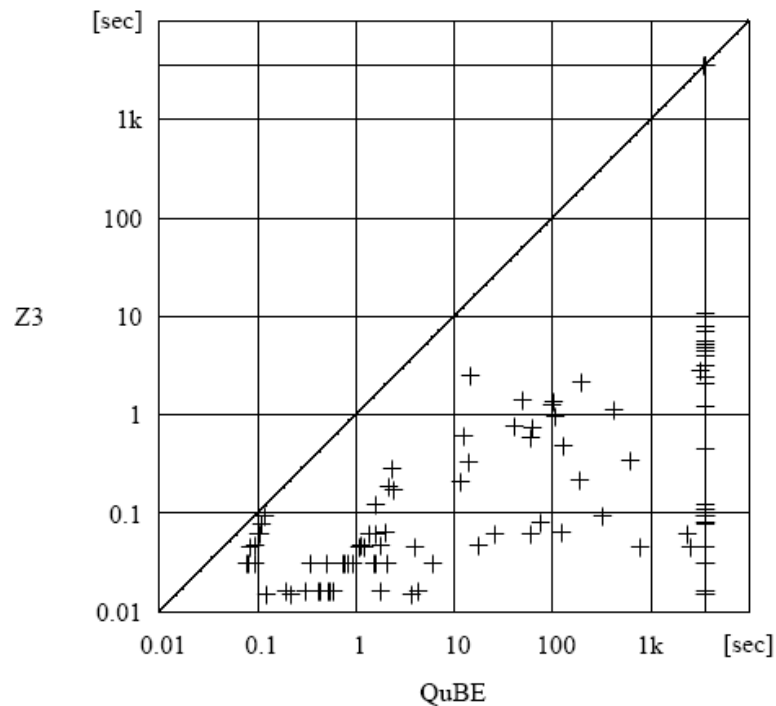
Hardware Fixpoint Checks.

Given: $I[x]$ and $T[x, x']$

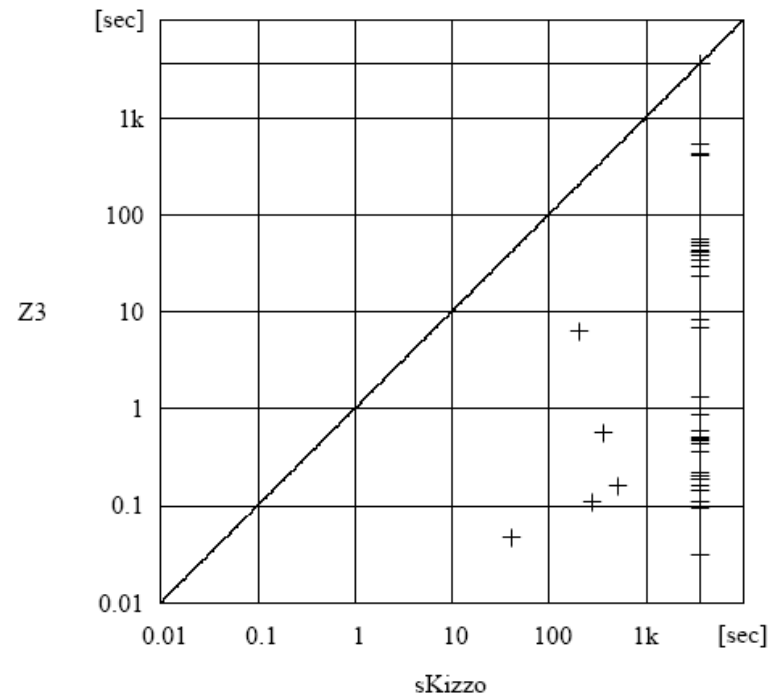
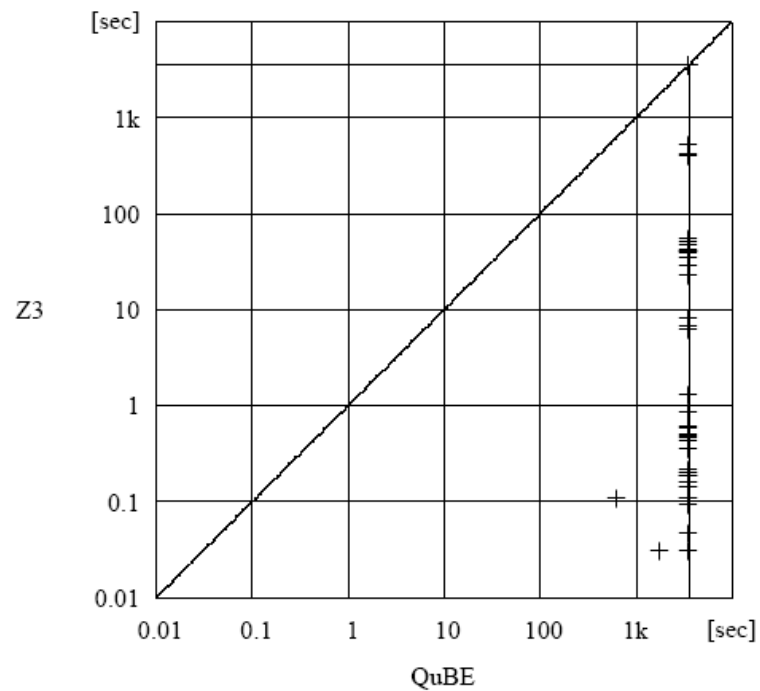
$$\forall x, x' . I[x] \wedge T^k[x, x'] \rightarrow \exists y, y' . I[y] \wedge T^{k-1}[y, y']$$

Ranking function synthesis.

Hardware Fixpoint Checks



Ranking Function Synthesis



Why is Z3 so fast in these benchmarks?

Z3 is using different engines:

rewriting, simplification, model checking, SAT, ...

Z3 is using a customized **strategy**.

We could do it because
we have access to the source code.

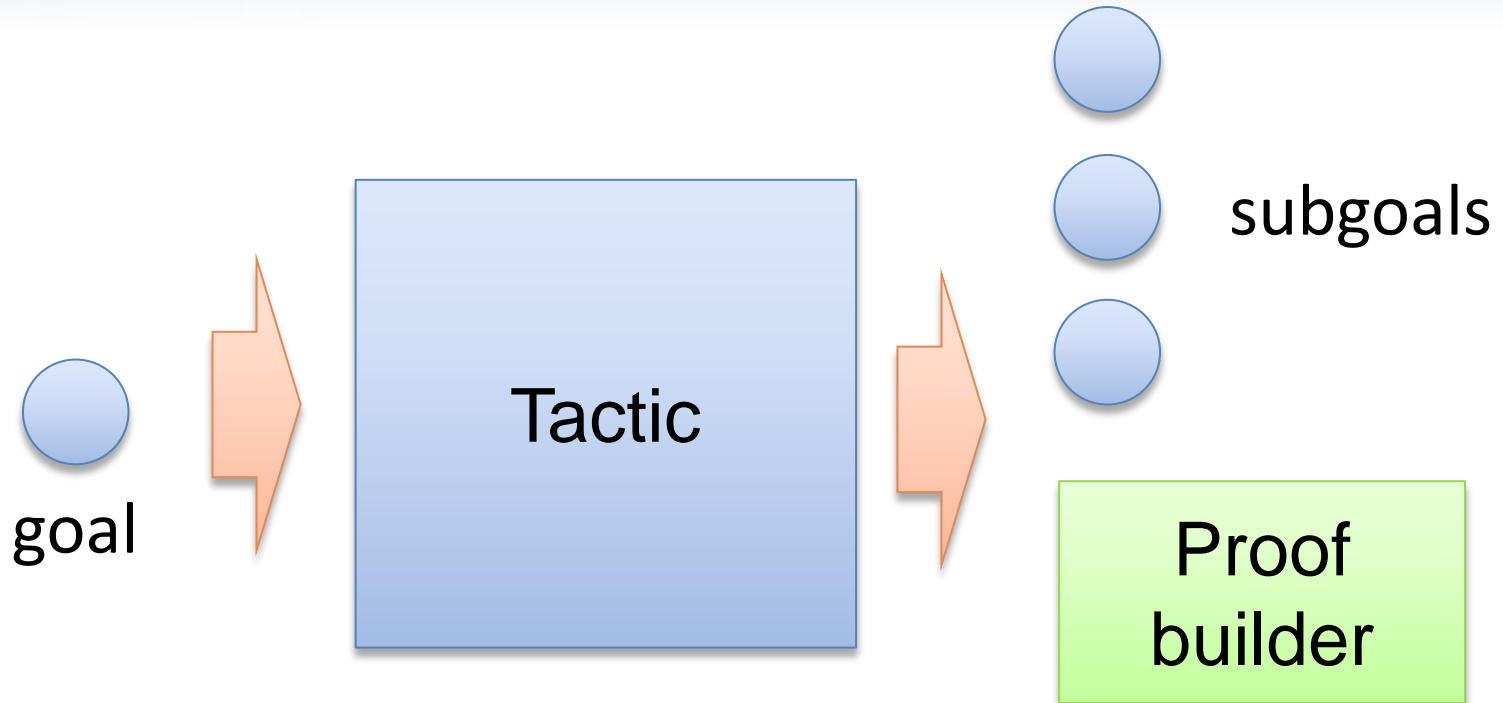
The "Message"

SMT solvers are collections of little engines.

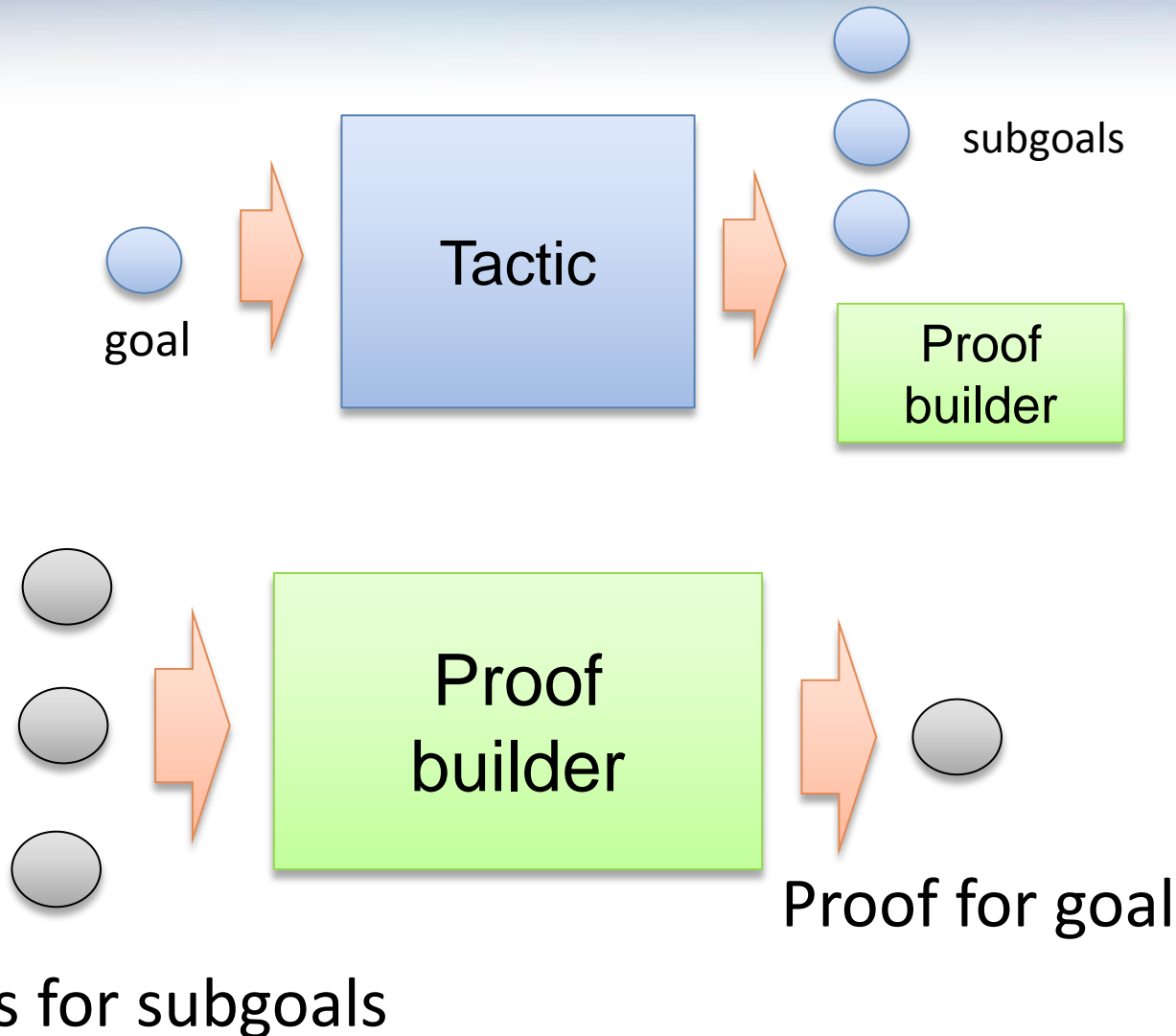
They should provide access to these engines.

Users should be able to define their own strategies.

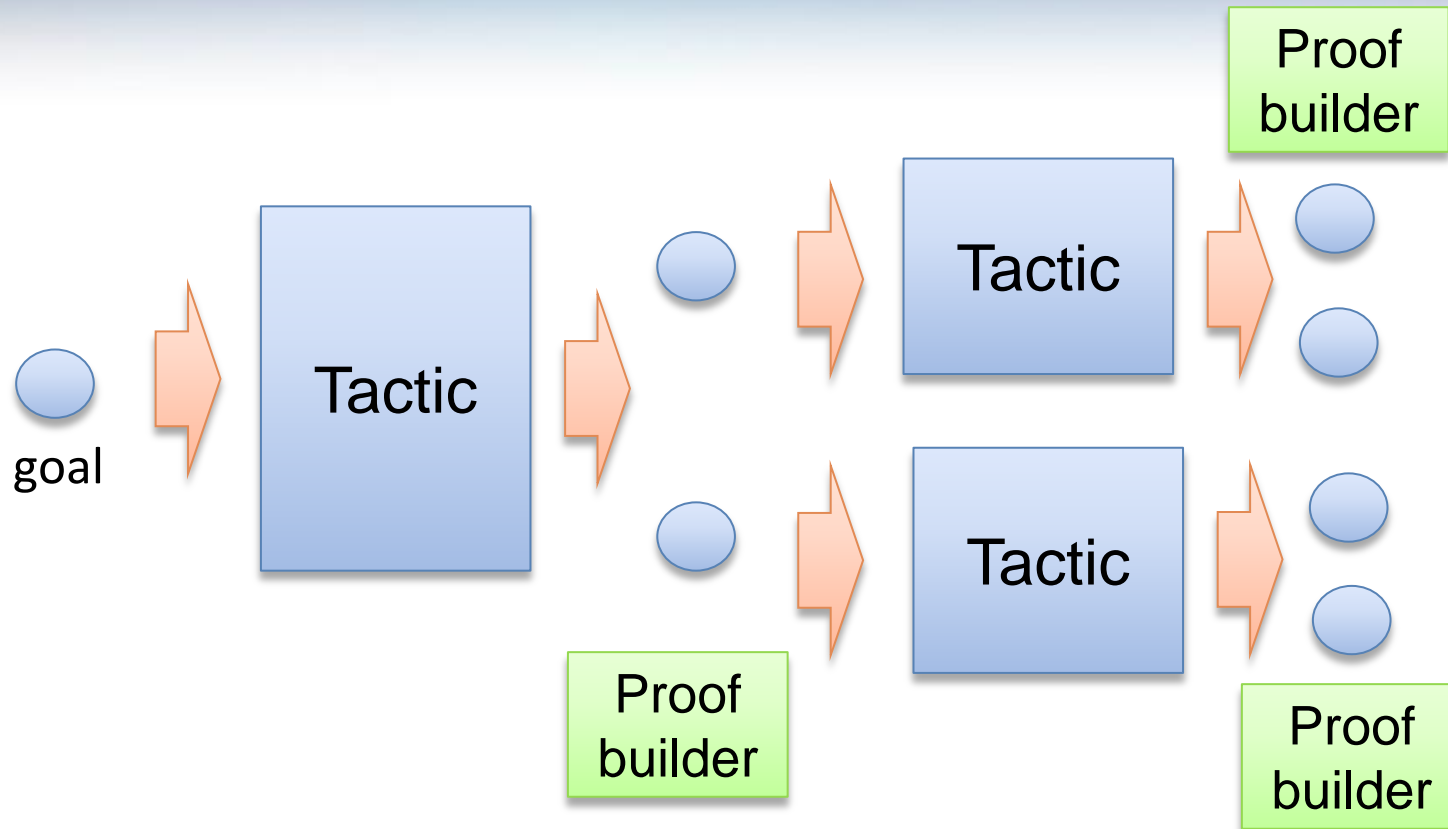
Main inspiration: LCF-approach



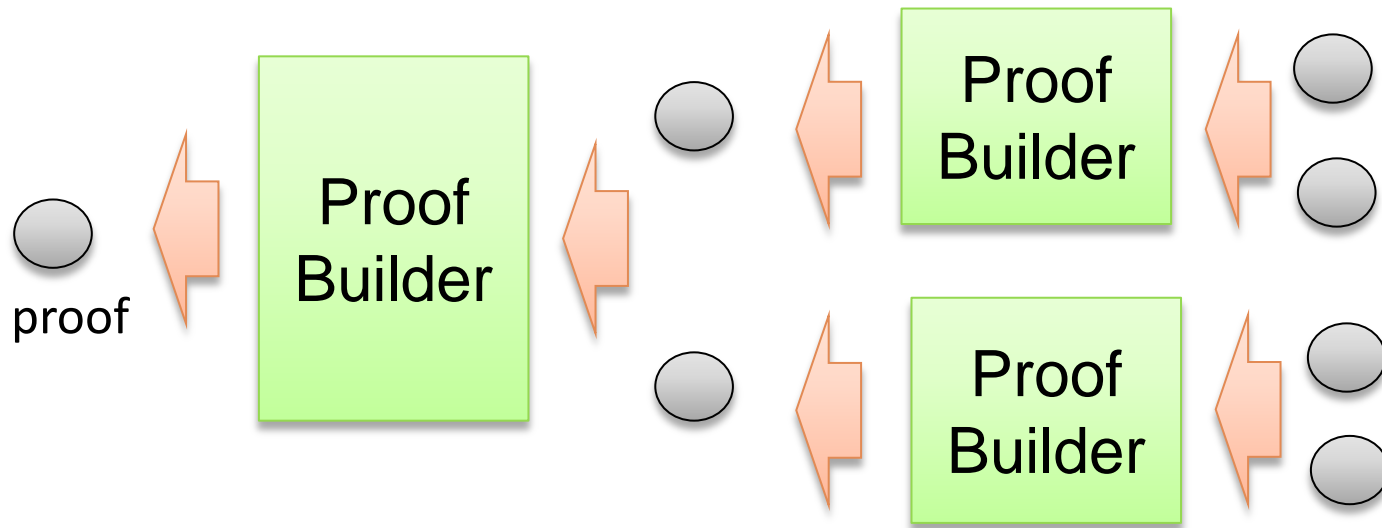
Main inspiration: LCF-approach



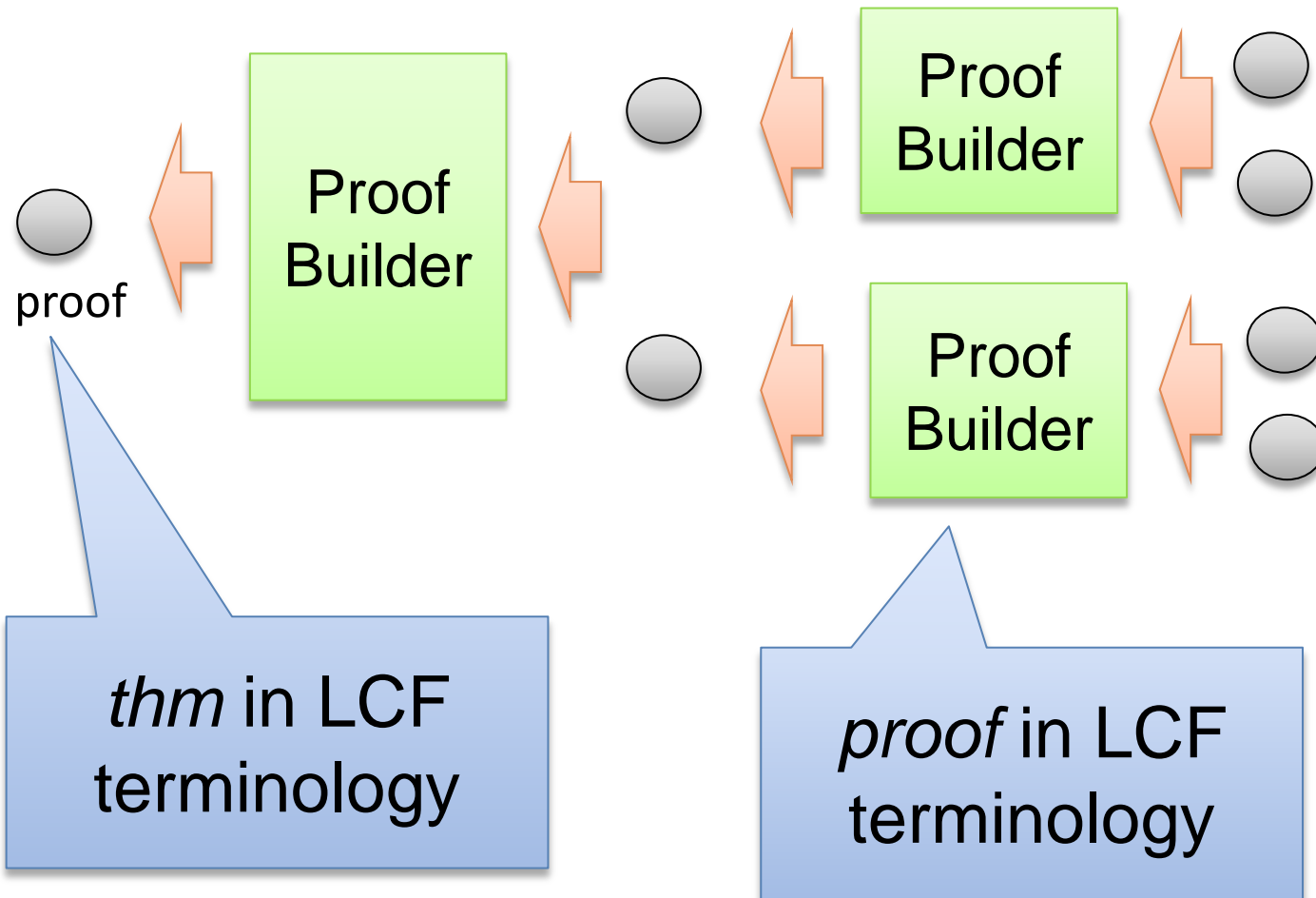
Main inspiration: LCF-approach



Main inspiration: LCF-approach





Main inspiration: LCF-approach



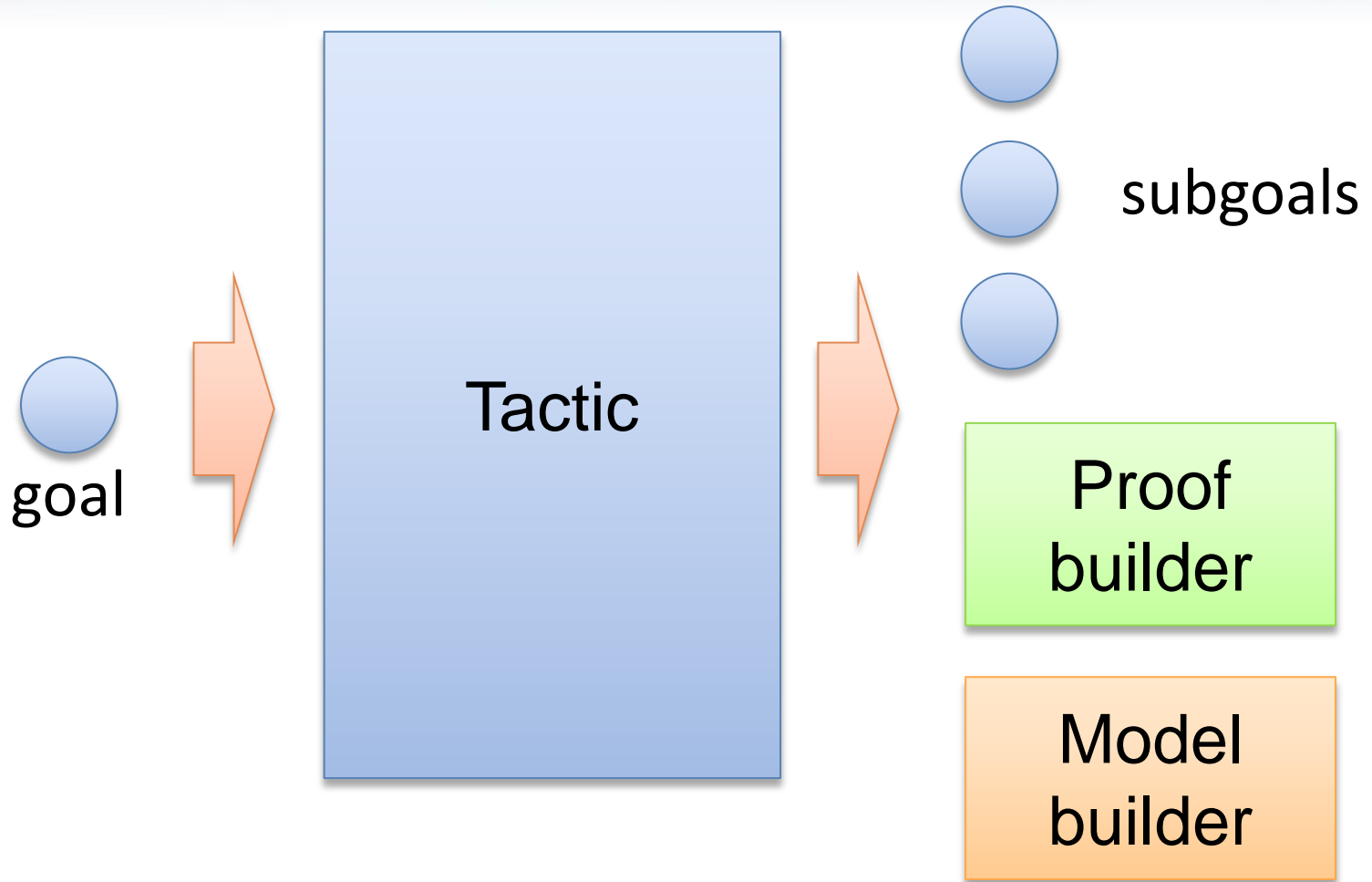
Tacticals aka Combinators

then( , ) = 

orelse( , ) = 

repeat() = 

SMT Tactic



SMT Tactic

goal = *formula sequence* \times *attribute sequence*

proofconv = *proof sequence* \rightarrow *proof*

modelconv = *model* \times *nat* \rightarrow *model*

trt = **sat** *model*

| **unsat** *proof*

| **unknown** *goal sequence* \times *modelconv* \times *proofconv*

| **fail**

tactic = *goal* \rightarrow *trt*

SMT Tactic

goal = *formula sequence* \times *attribute sequence*

proofconv = *proof sequence* \rightarrow *proof*

modelconv = *model* \times *nat* \rightarrow *model*

trt = **sat** *model*

| **unsat** *proof*

| **unknown** *goal sequence* \times *modelconv* \times *proofconv*

| **fail**

tactic = *goal* \rightarrow *trt*

end-game tactics:

never return unknown(sb, mc, pc)

SMT Tactic

goal = *formula sequence* \times *attribute sequence*

proofconv = *proof sequence* \rightarrow *proof*

modelconv = *model* \times *nat* \rightarrow *model*

trt = **sat** *model*

| **unsat** *proof*

| **unknown** *goal sequence* \times *modelconv* \times *proofconv*

| **fail**

tactic = *goal* \rightarrow *trt*

non-branching tactics:

sb is a singleton in
`unknown(sb, mc, pc)`

Trivial goals

Empty goal [] is trivially satisfiable

False goal [..., false, ...] is trivially unsatisfiable

basic : tactic

SMT Tactic example

$[a = b + 1, (a < 0 \vee a > 0), b > 3]$



Tactic:
elim-vars



Proof
builder

$[(b + 1 < 0 \vee b + 1 > 0), b > 3]$

Model
builder

SMT Tactic example

$[a = b + 1, (a < 0 \vee a > 0), b > 3]$



Tactic:
elim-vars



$[(b + 1 < 0 \vee b + 1 > 0), b > 3]$

Proof
builder

$M, M(a) = M(b) + 1$



Model
builder



M

SMT Tactic example

$[a = b + 1, (a < 0 \vee a > 0), b > 3]$



Tactic:
split-or



Proof
builder

$[a = b + 1, a < 0, b > 3]$

$[a = b + 1, a > 0, b > 3]$

Model
builder

SMT Tactics

simplify

nnf

cnf

tseitin

lift-if

bitblast

gb

mts

propagate-bounds

propagate-values

split-ineqs

split-eqs

rewrite

p-cad

sat

solve-eqs

SMT Tacticals

$\text{then} : (\text{tactic} \times \text{tactic}) \rightarrow \text{tactic}$

$\text{then}(t_1, t_2)$ applies t_1 to the given goal and t_2 to every subgoal produced by t_1 .

$\text{then}^* : (\text{tactic} \times \text{tactic sequence}) \rightarrow \text{tactic}$

$\text{then}^*(t_1, [t_{2_1}, \dots, t_{2_n}])$ applies t_1 to the given goal, producing subgoals g_1, \dots, g_m .

If $n \neq m$, the tactic fails. Otherwise, it applies t_{2_i} to every goal g_i .

$\text{orelse} : (\text{tactic} \times \text{tactic}) \rightarrow \text{tactic}$

$\text{orelse}(t_1, t_2)$ first applies t_1 to the given goal, if it fails then returns the result of t_2 applied to the given goal.

$\text{par} : (\text{tactic} \times \text{tactic}) \rightarrow \text{tactic}$

$\text{par}(t_1, t_2)$ executes t_1 and t_2 in parallel.

SMT Tacticals

$$\text{then}(\text{skip}, t) = \text{then}(t, \text{skip}) = t$$

$$\text{orelse}(\text{fail}, t) = \text{orelse}(t, \text{fail}) = t$$

SMT Tacticals

$\text{repeat} : \text{tactic} \rightarrow \text{tactic}$

Keep applying the given tactic until no subgoal is modified by it.

$\text{repeatupto} : \text{tactic} \times \text{nat} \rightarrow \text{tactic}$

Keep applying the given tactic until no subgoal is modified by it, or the maximum number of iterations is reached.

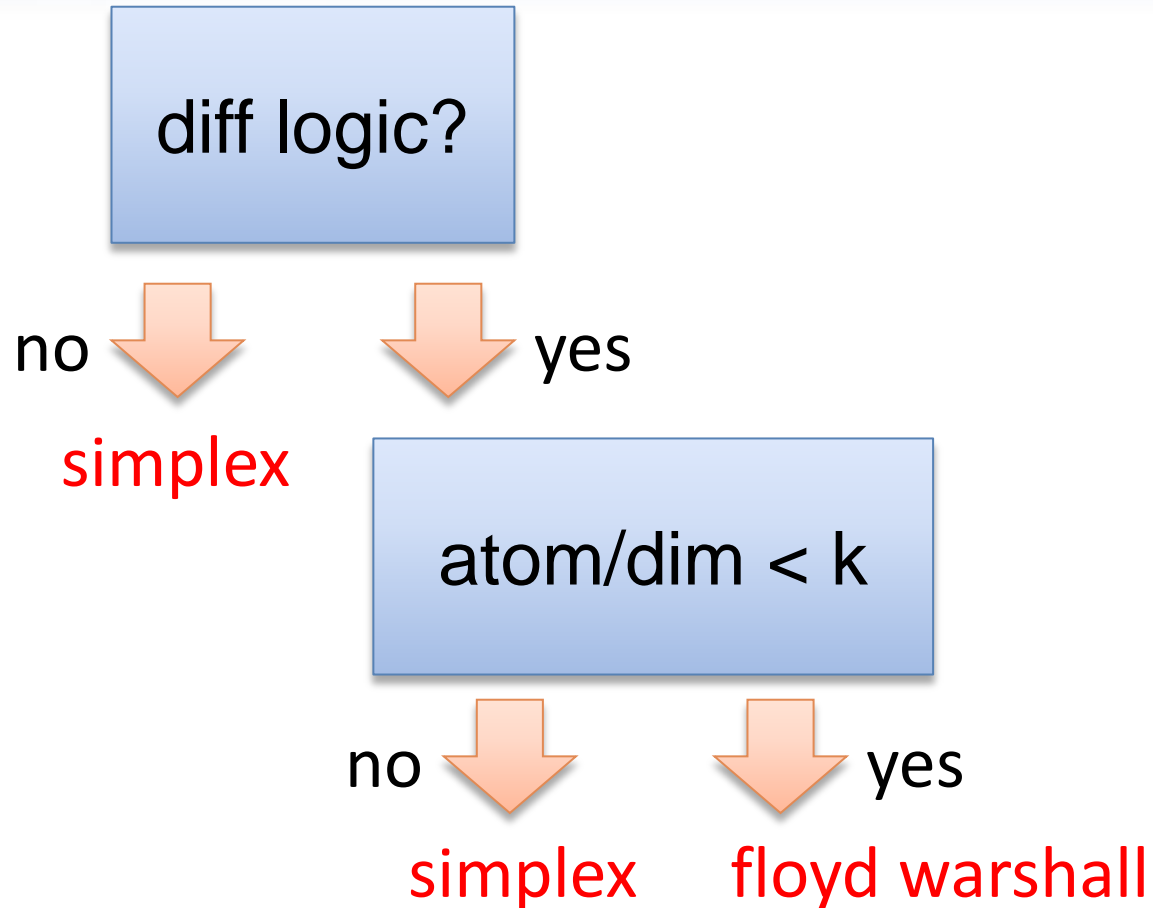
$\text{tryfor} : \text{tactic} \times \text{seconds} \rightarrow \text{tactic}$

$\text{tryfor}(t, k)$ returns the value computed by tactic t applied to the given goal if this value is computed within k seconds, otherwise it fails.

Feature / Measures

Probing structural features of formulas.

Feature / Measures: Yices Strategy



Feature / Measures: Yices Strategy

orelse(then(failif($\text{diff} \wedge \frac{\text{atom}}{\text{dim}} > k$), simplex), floydwarshall)

Fail if condition is not satisfied.
Otherwise, do nothing.

Feature / Measures: Examples

bw: Sum total bit-width of all rational coefficients of polynomials in case.

diff: True if the formula is in the difference logic fragment.

linear: True if all polynomials are linear.

dim: Number of arithmetic constants.

atoms: Number of atoms.

degree: Maximal total multivariate degree of polynomials.

size: Total formula size.

Tacticals: syntax sugar

$\text{if}(c, t_1, t_2) = \text{orelse}(\text{then}(\text{failif}(\neg c), t_1), t_2)$
 $\text{when}(c, t) = \text{if}(c, t, \text{skip})$

Under/Over-Approximations

Under-approximation

unsat answers cannot be trusted

Over-approximation

sat answers cannot be trusted

Under/Over-Approximations

Under-approximation
model finders

Over-approximation
proof finders

Under/Over-Approximations

Under-approximation

$$S \rightarrow S \cup S'$$

Over-approximation

$$S \rightarrow S \setminus S'$$

Under/Over-Approximations

Under-approximation

Example: QF_NIA model finders
add bounds to unbounded variables (and blast)

Over-approximation

Example: Boolean abstraction

Under/Over-Approximations

Combining under and over is bad!
sat and unsat answers cannot be trusted.

Tracking: under/over-approximations

In principle, proof and model converters can check if the resultant models and proofs are valid.

Tracking: under/over-approximations

In principle, proof and model converters can check if the resultant models and proofs are valid.

Problem: if it fails what do we do?

Tracking: under/over-approximations

In principle, proof and model converters can check if the resultant models and proofs are valid.

Problem: if it fails what do we do?

We want to write tactics that can check whether a goal is the result of an abstraction or not.

Tracking: under/over-approximations

Solution

Associate an **precision attribute** to each goal.

Goal Attributes

Store extra logical information

Examples:

precision markers

goal depth

polynomial factorizations

Decision Engines as Tacticals

AP-CAD (tactic) = tactic

Decision Engines as Tacticals

```
then(preprocess, smt(finalcheck))
```

Strategy: Example

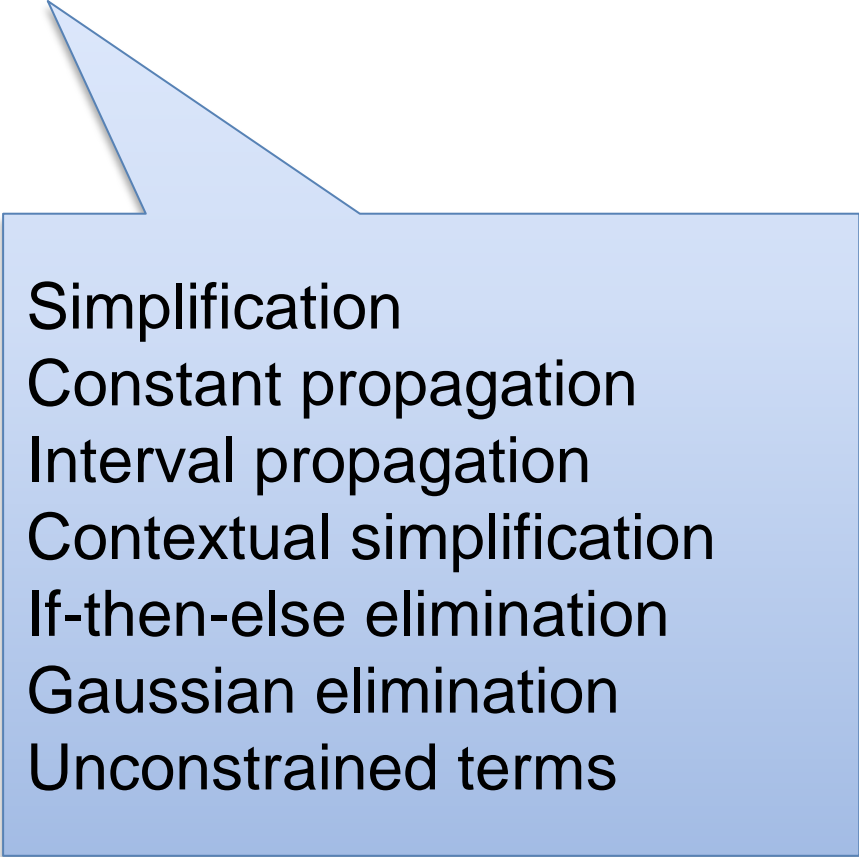
```
then(then(simplify, gaussian), orelse(modelfinder, smt(apcad(icp))))
```

RAHD Calculemus Strategy

	dim	deg	div	calc-0	calc-1	calc-2	qepcad-b	redlog/rlqe	redlog/rlcad
P0	5	4	N	.91	1.59	1.7	416.45*	40.4	-
P1	6	4	N	1.69	3.08	3.42	.*	-	-
P2	5	4	N	1.34	2.41	2.62	.*	-	-
P3	5	4	N	1.52	2.56	2.75	.*	-	-
P4	5	4	N	1.14	2.02	2.16	.*	-	-
P5	14	2	N	.25	.26	.27	.*	97.4	-
P6	11	5	N	147.4	.07	.06	.*	<.01	<.01
P7	8	2	N	.05	<.01	<.01	.08	<.01	<.01
P8	7	32	N	4.5	.1	<.01	8.38	<.01	-
P9	7	16	N	4.51	.15	<.01	.29	.01	6.7
P10	7	12	N	100.74	20.76	8.85	.*	-	-
P11	6	2	Y	1.6	.5	.53	.01	.01	.05
P12	5	3	N	.78	.3	.36	.02	.01	.07
P13	4	10	N	3.83	3.95	4.02	.*	-	-
P14	2	2	N	4.55	1.67	.07	.01	-	-
P15	4	3	Y	.177	.2	.12	.01	<.01	<.01
P16	4	2	N	9.99	2.17	2.1	.02	<.01	<.01
P17	4	2	N	.62	.59	.65	.28	.02	.61
P18	4	2	N	1.25	1.28	1.27	.01	<.01	<.01
P19	3	6	Y	3.34	1.72	2.08	.02	.01	.7
P20	3	4	N	1.18	.65	.65	.01	<.01	.3
P21	3	2	N	.02	.03	<.01	.02	.01	.1
P22	2	4	N	<.01	<.01	<.01	.01	<.01	<.01
P23	2	2	Y	<.01	<.01	<.01	<.01	<.01	<.01

Z3 QF_LIA Strategy

```
then(preamble, orelse(mf, pb, bounded, smt))
```



Simplification
Constant propagation
Interval propagation
Contextual simplification
If-then-else elimination
Gaussian elimination
Unconstrained terms

Challenge: small step configuration

proof procedure as a transition system

Abstract DPLL, DPLL(T), Abstract GB, cutsat, ...

UnitPropagate :

$$M \parallel F, C \vee l \implies M \parallel l \parallel F, C \vee l \quad \text{if} \quad \begin{cases} M \models \neg C \\ l \text{ is undefined in } M \end{cases}$$

PureLiteral :

$$M \parallel F \implies M \parallel l \parallel F \quad \text{if} \quad \begin{cases} l \text{ occurs in some clause of } F \\ \neg l \text{ occurs in no clause of } F \\ l \text{ is undefined in } M \end{cases}$$

Decide :

$$M \parallel F \implies M \parallel l^d \parallel F \quad \text{if} \quad \begin{cases} l \text{ or } \neg l \text{ occurs in a clause of } F \\ l \text{ is undefined in } M \end{cases}$$

Fail :

$$M \parallel F, C \implies \text{FailState} \quad \text{if} \quad \begin{cases} M \models \neg C \\ M \text{ contains no decision literals} \end{cases}$$

Backtrack :

$$M \parallel l^d \parallel N \parallel F, C \implies M \parallel \neg l \parallel F, C \quad \text{if} \quad \begin{cases} M \parallel l^d \parallel N \models \neg C \\ N \text{ contains no decision literals} \end{cases}$$

Challenge: small step configuration

proof procedure as a transition system

Abstract DPLL, DPLL(T), Abstract GB, cutsat, ...

Challenge:

Efficient strategic control

Decide : $M \parallel F \Rightarrow M \parallel F$ if $\begin{cases} l \text{ or } \neg l \text{ occurs in a clause of } F \\ l \text{ is undefined in } M \end{cases}$

Fail : $M \parallel F, C \Rightarrow \text{FailState}$ if $\begin{cases} M \models \neg C \\ M \text{ contains no decision literals} \end{cases}$

Backtrack :

$M \parallel N \parallel F, C \Rightarrow M \parallel \neg l \parallel F, C$ if $\begin{cases} M \parallel N \models \neg C \\ N \text{ contains no decision literals} \end{cases}$

Conclusion

Different domains need different strategies.

We must expose the little engines in SMT solvers.

Interaction between different engines is a must.

Tactic and Tacticals: **big step approach**.

More transparency.